

The **Clarke** (aka. **VCG**) mechanism [Clarke 71]

- The Clarke mechanism chooses some outcome o that maximizes $\sum_i v_i(\theta_i', o)$
 - θ_i' = the type that i reports
- To determine the payment that agent j must make:
 - Pretend j does not exist, and choose o_{-j} that maximizes $\sum_{i \neq j} v_i(\theta_i', o_{-j})$
 - j pays $\sum_{i \neq j} v_i(\theta_i', o_{-j}) - \sum_{i \neq j} v_i(\theta_i', o)$

The Clarke mechanism is strategy-proof

- Total utility for agent j is
$$v_j(\theta_j, o) + \sum_{i \neq j} v_i(\theta_i', o) - \sum_{i \neq j} v_i(\theta_i', o_{-j})$$
- But agent j cannot affect the choice of o_{-j}
- Hence, j can focus on maximizing $v_j(\theta_j, o) + \sum_{i \neq j} v_i(\theta_i', o)$
 - The other agents' types θ_i' are out of j 's control
 - Agent j 's true type θ_j is also fixed
 - Agent j will try to report a type θ_j' so that o maximizes $v_j(\theta_j', o) + \sum_{i \neq j} v_i(\theta_i', o)$
- But mechanism chooses o to maximize $\sum_i v_i(\theta_i', o)$
- Hence, if $\theta_j' = \theta_j$, o maximizes $\sum_i v_i(\theta_i', o) = v_j(\theta_j', o) + \sum_{i \neq j} v_i(\theta_i', o)$. Therefore, if $\theta_j' = \theta_j$, o maximizes $v_j(\theta_j, o) + \sum_{i \neq j} v_i(\theta_i', o)$.