Memory Model

- **For this course: Assume *Uniform Access Time***
  - All elements in an array accessible with same time cost
  - Reality is somewhat different

- **Memory Hierarchy (in order of decreasing speed)**
  - Registers
  - On (cpu) chip cache memory
  - Off chip cache memory
  - Main memory
  - Virtual memory (automatically managed use of disk)
  - Explicit disk I/O

- **All but last managed by system**
  - Need to be aware, but can do little to manipulate others directly
  - Promote locality?
Cost of Disk I/O

- Disk access 10,000 to 100,000 times slower than memory access
  - Do almost anything (almost!) in terms of computation to avoid an extra disk access
  - Performance penalty is huge

- B-Trees designed to be used with disk storage
  - Typically used with database application
  - Many different variations
  - Will present basic ideas here

- Want broad, not deep trees
  - Even log N disk accesses can be too many
External Methods

- Disk use requires special consideration
  - Timing Considerations (already mentioned)
  - Writing pointers to disk?
  - What do values mean when read in at a different time/different machine?

- General Properties of B-Trees
  - All Leaves Have Same Depth
  - Degree of Node > 2
  - (maybe hundreds or thousands)

- Not a Binary Tree, but is a Search Tree
  - There are many implementations...
  - Will use examples with artificially small numbers to illustrate
Rules of B-Trees

**Rules**

1. Every node (except root) has at least MINIMUM entries
2. The MAXIMUM number of node entries is 2*MINIMUM
3. The entries of each B-tree are stored, sorted
4. The number of sub-trees below a non-leaf node is one greater than the number of node entries
5. For non leaves:
   - Entry at index k is greater than all entries in sub-tree k
   - Entry at index k is less than all entries in sub-tree k+1
6. Every leaf in a B-tree has the same depth
Example B Tree (MAX = 2)

[6]

*  *

*  *

*  *

[2  4]  [9]

*  *  *

*  *

*  *

*  *

[1]  [3]  [5]  [7  8]  [10]
Search in B-Tree

- Every Child is Also the Root of a Smaller B-Tree
- Possible internal node implementation
  ```java
class BTNode {
    int myDataCount;
    int myChildCount;
    KeyType[] myKeys[MAX+1];
    BTNode[] myChild[MAX+2]
}
```
- Search:
  ```java
  boolean isInBTree(BTNode t, KeyType key);
  1. Search through myKeys until myKeys[k] >= key
  2. If t.myData[k] == key, return true
  3. If isLeaf(t) return false
  4. return isInBtree(t.myChild[k])
  ```
Find Example

Example Find in B-Tree (MAX = 2)

Finding 10
B-Tree Insertion

- **Insertion Gets a Little Messy**
  - Insertion may cause rule violation
  - “Loose” Insertion (leave extra space) (+1)
  - Fixing Excess Entries

- **Insert Fix**
  - Split
  - Move up middle
  - Height gained only at root

- **Look at some examples**
Insertion Fix

- (MAX = 4) Fixing Child with Excess Entry

Before: [9 28]

[3 4] [13 16 19 22 25] [33 40]

[2 3] [4 5] [7 8] [11 12] [14 15] [17 18] [20 21] [23 24] [26 27] [31 32] [34 35] [50 51]

After: [9 19 28]

[3 4] [13 16] [22 25] [33 40]

[2 3] [4 5] [7 8] [11 12] [14 15] [17 18] [20 21] [23 24] [26 27] [31 32] [34 35] [50 51]
Insertion Fix

(MAX=2) Another Fix

[6 17] BEFORE
[4] [12] [18 19 22]

[6 17 19] STEP 1
[4] [12] [18] [22]

STEP 2

[ ]

[6 17 19]
[4] [12] [18] [22]

[17] AFTER

[6] [19]

CompSci 100E
B-Tree Removal

- **Remove**
  - Loose Remove
  - If rules violated: Fix
    - Borrow (rotation)
    - Join
- **Examples left to the “reader”**
B-Trees

- **Many variations**
  - Leaf node often different from internal node
  - Only leaf nodes carry all data (internal nodes: keys only)
  - Examples didn’t distinguish keys from data
  - Design to have nodes fit disk block

- **The Big Picture**
  - Details can be worked out
  - Can do a lot of computation to avoid a disk access